## Botnet Population and Intelligence Gathering Techniques

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# BlackHat DC Meeting 2008



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## Introductions



The Spacious Georgia Tech Campus based on joint work with:

- UCF CS: Cliff Zou
- GaTech CS: Jason Trost, Wenke Lee
- ISC: Paul Vixie
- IOActive: Dan Kaminski
- Thanks: Nicholas Bourbaki



## Outline

- Motivation: Infer victim populations with limited probes
- IPID overview
- BIND Cache Overview
- Challenges in Modeling
- Solutions
- Further challenges
- Data needs: finding honest open recursives
- Cautions and conclusions



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#### **Basic Botnet Facts**

- Most bot malware will utilize domain names so the bot master can move around and the bots can still find him.
- Many types of bot malware use multiple staged downloads.
- Many bot masters are just starting to understand how to get their bots to egress from corporate networks.
- Alot of bot malware is shockingly easy to use

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#### **Botnet Basics: Rats**



David Dagon & Chris Davis Botnet Population Estimation

#### **Botnet Basics: Rats**

🗏 Poison Ivy		
	Advanced [2]	Size: 6.20 KiB
Profiles	Process Mutex: JVoqA.14 ()	
p Connection	Inject server into the default browser     Persistence     Inject into a running process	
🤹 Install	Process: mannagr.exe Key logger Format:	
<del>%</del> Advanced	<ul> <li>G PE</li> <li>File Alignment (bytes): 512</li></ul>	
Build	<ul> <li>☞ Binary</li> <li>© CArray</li> <li>© Delphi Array</li> <li>© Python Array</li> </ul>	
	Cancel	Next ⇒
Version 2.3.1 Nr	of Ports: 0 Nr. of Plugins: 2 Nr. of Connections: 0	



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#### **Botnet Basics: Rats**





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#### **Basic Botnet Facts**

- Not Your Mom's IRC Botnet anymore
- IRC Botnets are on the decline. Remote Victim Enumeration is becoming harder
- How do we understand the size and scope of a botnet when we have a limited view?



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- Each IP datagram header has an ID field, which is used when reassembling fragmented datagrams.
- If no fragmentation takes place, the ID field is basically unused, but operating systems still have to calculate its value for each packet.
- Some operating systems increment the value by a constant for each datagram.
- Operating systems that increment by one:
  - Windows (All Versions)
  - FreeBSD
  - Some Linux Variants (2.2 and Earlier)
  - Many other devices like print servers, webcams, etc...



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An example of a quiet server: cdavis\$ hping2 -i 1 -c 5 -S -p 80 XX.YY.ZZ.86 len=46 ip=XX.YY.ZZ.86 ttl=52 id=25542 sport=80 flags=SA seg=0 win=8192 rtt=42.2 ms len=46 ip=XX.YY.ZZ.86 ttl=52 id=25543 sport=80 flags=SA seg=1 win=8192 rtt=48.6 ms len=46 ip=XX.YY.ZZ.86 ttl=52 id=25544 sport=80 flags=SA seg=2 win=8192 rtt=48.1 ms len=46 ip=XX.YY.ZZ.86 ttl=52 id=25545 sport=80 flags=SA seg=3 win=8192 rtt=43.9 ms

len=46 ip=XX.YY.ZZ.86 ttl=52 id=25546
sport=80 flags=SA seq=4 win=8192 rtt=42.1 ms

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## Motivation

- 80% of spam sent via zombies [St.Sauver 2005]; now 90+% [St.Sauver 2007]
- Volume of phish/malware complaints to ISPs is staggering
  - Need to prioritize
- So-called IP-reputation is often merely CIDR-Reputation
  - DHCP auto-incrementing spam bots, and general lease churn mitigates towards classful scoring, or based on whois OrgName or ASN, etc.
  - Need to remotely assess risk of networks roughly (CIDR) 2 without relying on remote sensors.
- Motivating guestion: Can we estimate victim populations using simple DNS metrics?



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#### • Epidemiological Studies via DNS Cache:





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• Epidemiological Studies via DNS Cache:



#### • Intuitive Difference in Relative Cache Rates





## **Conception Application of DNS Cache Snooping**

• Probing Caching Servers for Same Domain





#### • Caching Inherently Hides Lookups



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## Solution: Boundary Estimates

#### Assumptions

- Property 1: Bot queries are independent
- Property 2: DNS Cache queues follow a Poisson distribution with the arrival of uncached phases at rate λ
  - Note: λ is the "birth process", or arrival rate-the number of events/arrivals per time epoch.
- Are these properties correct?



## Independence of Bot Queries

- Two events X<sub>i</sub> and X<sub>i</sub>, are independent if
  - $P(X_i X_i) = P(X_i)P(X_i)$
  - Given the property that P(B|A) = P(BA)/P(A), then to show  $X_i$  and  $X_i$  are independent, we need to show  $P(X_i|X_i) = P(X_i)$
- In the general case, bot victims are randomly selected from potential victims.
- Absent synchronized behavior, one victim's infection-phase DNS resolution is independent of any others.
- Example: two victims must visit a webpage to become infected; on a domain TTL-scale, this browsing is independent
- Thus, proptery 1 holds in the general case

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## Bot DNS Resolution Follows Poisson Distribution

• Does Property 2 hold? Consider:

#### Intuitive View of DNS Cache Time-outs



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## Bot DNS Resolution Follows Poisson Distribution

- The arrival of victims in a queue is trivially modeled as a poisson process
  - This is true of telephony networks, packet networks
  - ...and its generally true of origination from *large* populations of independent actors
- (For some values of large) botnets are large population systems.
- OK, so keep in mind: botnet recruitment that triggers a DNS lookup is a poisson process. We use this point shortly...
- Our current problem: We can only measure cache idle periods however. Are these poisson processes?



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#### Motivation

#### Poisson Processes Definitions

- What's a Poisson process? There are three definitions:
  - One arrival occurs in the infinitesimal time dt
  - An interval t has a distribution of arrivals following  $P(\lambda t)$
  - The interarrival times are independent with exponential distribution.  $P\{interarrival > t\} = e^{-\lambda t}$
- Say, that third definition sure looks like a DNS cache line's idle periods!
- Textbooks then tell used:  $\hat{N}_{u,l} = \hat{\lambda}_{u,l}/\lambda$ . (There are simple models for deriving populations from arrival rates.)
- Bad joke opportunity: DNS poisoning also relies on poisson processes



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#### **More Problems**

#### • There are hazards in sampling

- Hidden masters
- Load balancers using independent caches
- Policy barriers
- Mandatory
  - Obtain permission and follow RFC 1262 (DNS probes are the spam)
  - Throttle request rates to respect server load balancing (or corrupt data); e.g., 4.2.2.2 throttles non-customers
  - Select small set of suspect domains
- All of these corrupt data collection.
  - (Solutions omitted for space)

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#### Sampling is Blind to DNS Architecture





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## Sample Application

Study of botnet in Single ISP DNS Cache





#### Demonstration

#### Plot of output for tracking one botnet (animation may follow)





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#### Issue: How to Locate Open Recursives?

- Probing open recursives for domain cache times requires a list of open resolvers.
  - We could just ... scan IPv4 for such hosts
- However, simple queries don't tell us the whole story of the open recursives needed for this task
- We must separate those that are open recursive from those that are open forwarding
- Further, some open resolvers (both full and forwarding) are DNS monetization engines, and don't answer iterative queries truthfully
  - DNS monetization resolvers may not uses caches
  - We wish to identify them, so we can exclude them



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Motivation

## One Approach to Recursive/Forwarding Enumeration



## Study Methodology



- Unique label queried to all IPv4
- SOA wildcard for parent zone
- Script used to return srcIP of requester
- Logging at NS yields open recursive and recursive forwarding hosts
- Further analysis enumerates "interesting" resolvers



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## Methodology (cont'd)

#### Phase1

- If response given...
- Exclude authority open resolvers
- fpdns taken of answering host
- Perform http request of host
- Phase2
  - Pick 600K open resolvers
  - Ask them repeatedly to resolve phishable domains
  - Note which ones gave incorrect answers
  - If "incorrect", http request to the answered IP



Motivation

## Open Recursion: Comparison of /16s, in IPv4



Open Recursive Hosts in /16 CIDRs

Motivation

#### Open Recursion: Comparison of /16s, in IPv4



Motivation

#### Open Recursion: Putative GNU libc /16s



Open Recursive Hosts in /16 CIDRs



Motivation

## Open Recursion: Putative GNU libc /16s



- gnu libc logic of AAAA? → A? queries.
- Other heuristics: Windows DNS servers answered authoritatively for queries for
  - 1.in-addr.arpa,
- Someone needs to update fpdns (2005)
- Other "harmless" explanations for open recursion can be considered, and accepted or discarded



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Motivation

## Open Recursion: Histogram of Queries to NS





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#### Analysis: What DNS Server is Running?

#### • HTTP server string fetched from open recursive hosts

- $\bullet~\sim$  20% RomPager, Nucleus, misc. known devices
- $\sim$  80% No answer
- Thus, designed study groups:
  - Randomly selected open recursive resolvers
  - Intersection of open recursives and visitors to Google's authority server
  - Intersection of open recursives and Storm victims



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## Filtering Out "Non-Spec" DNS Servers

- Methodology:
  - Selected 200K random open recs, 200K open recs contacting Google authority servers, 200K overlap storm
  - Repeatedly queried for "phishable"; 15 min window; 220M probes total over 4 days
  - Diurnal pattern noted (unusual for DNS servers)
  - Approx. 310K-330K resolvers answer; 460K out of 600K total answered
- 2.4% were technically "incorrect" (extrapolates to 291,500K hosts)
- 0.4% were malicious (extrapolates to 68K hosts; 36K measured so far in subsequent full IPv4 sweeps)



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## Filtering Out "Non-Spec" DNS Servers

Created database of "proxied" webpages

- Porn, advertising, and proxied pages(!)
- ho  $\sim$  20% proxied/rewrote google.com (demo)
- $\bullet~\sim$  11% proxied a chinese search page
- $\bullet~\sim$  26% proxied a comcast user login
- Methodology reported in www.isoc.org/isoc/conferences/ndss/08
- In short, we need to remove these hosts from our open recursive pool



#### Filtering out "Non-Spec" DNS: Why?

#### Baaaad DNS (and therefore bad cache timing data):





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## Conclusions

- DNS cache inspection requires careful analysis
- Merely probing DNS caches alone does not reveal victim information
- A model (with safe assumptions) is needed to overcome noise created by variable DNS architecture, events, etc.



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- Merely probing DNS caches alone does not reveal victim information
- A model (with safe assumptions) is needed to overcome noise created by variable DNS architecture, events, etc.

#### Notify, Ask and Coordinate

- Uncoordinated DNS probes pollute IDS logs, generate e-mail complaints
- Use RFC 1262, and common courtesy
- Don't bother checking mil or gov prefixes



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